

# Single shot q-thermo, VL 2

Course homepage: [www.mpinwuerzburg.net/lecture2](http://www.mpinwuerzburg.net/lecture2)

No lecture/tutorial/exercise sheet in the week  
of Nov. 3-7!

Tutorial: Tuesday (Di.) 14:15 - 15:45,  
Philosophenweg 12, room 060

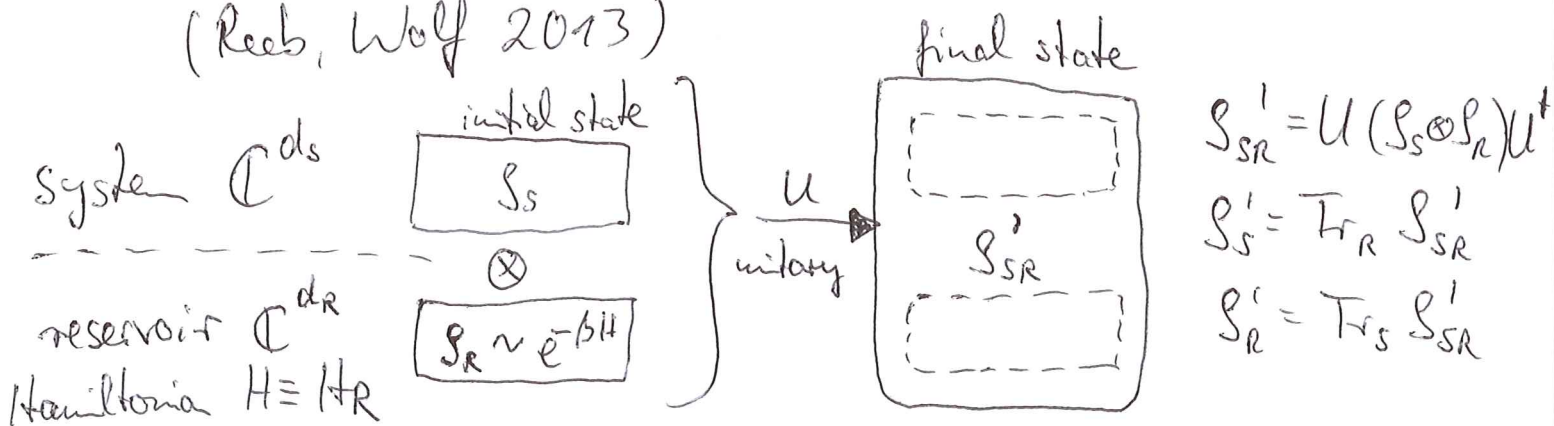
Credit points for 60% of homework.

Goal: << work than theor. Phys. exercises  
groups of 2 or 3 people.

any further organisational questions?

## 1. Landauer's Principle: rigorous derivations

### 1.1. A quantum-information-theoretic version (Reeb, Wolf 2013)



$\Delta S_S := S(S_S) - S(S'_S)$  entropy reduction in system

$\Delta S_R := S(S_R) - S(S'_R)$  " " " reservoir

①  $\Delta Q := \text{tr}[H S'_R] - \text{tr}[H S_R]$  heat increase in reservoir.

Theorem:  $\Delta Q \geq k_B T \Delta S_S$

$\beta = 1/(k_B T)$

In more detail

$$\Delta Q = k_B T \left( \Delta S_S + \underbrace{I(S':R')}_{\substack{\text{mutual information} \\ \geq 0}} + \underbrace{D(S_R' \| S_R)}_{\substack{\text{relative entropy} \\ \geq 0}} \right)$$

Remarks:

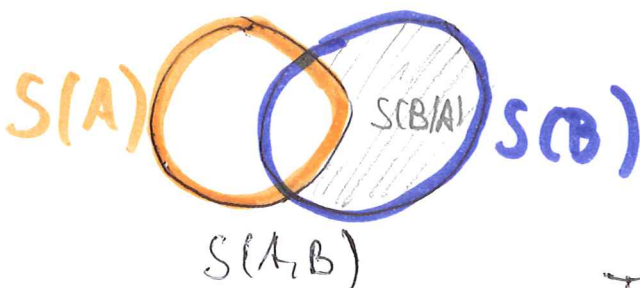
• This is work extraction on average (exp. value of  $W = W_R$ ; to be determined in many runs of the experiment).

"Single-shot" approach will be: run only once, and aim at prob  $\geq 1 - \epsilon$  for dissipation of energy  $\leq E$ .

•  $I(S':R')$  is a measure of correlations in the final state:

$$\begin{aligned} I(A:B) &= S(A) + S(B) - S(AB) \\ &\equiv S(S_A) + S(S_B) - S(S_{AB}) \\ &= S(B) - S(B|A) \geq 0 \end{aligned}$$

where  $S(B|A) := S(A,B) - S(A)$  is the conditional entropy: remaining uncertainty about B if we learn about A



$\Rightarrow I(A:B) =$  how much learning A helps us to learn about B.

$$I(A:B) = I(B:A)$$

②

Theorem says: correlations are costly!

Ex.:  $\rho_{AB} = \rho_A \otimes \rho_B$  (uncorrelated)  $\Rightarrow S(AB) = S(A) + S(B)$

$$\Rightarrow I(A:B) = 0.$$

$$\rho_{AB} = \frac{1}{2} |00\rangle\langle 00| + \frac{1}{2} |11\rangle\langle 11| \quad (\text{classical correlation})$$

$$\Rightarrow \rho_A = \rho_B = \frac{1}{2} \mathbb{1} \Rightarrow S(AB) = S(A) = S(B) = \log 2$$

$$\Rightarrow I(A:B) = \log 2 \quad [=1 \text{ if } \log = \log_2]$$

"one bit of mutual information".

Proof: First we show the "second law lemma"

$$\underbrace{(-\Delta S_S)}_{\text{entropy increase in system}} + \underbrace{(-\Delta S_R)}_{\text{entropy increase in reservoir}} = I(S':R') \geq 0.$$

$$\begin{aligned} -\Delta S_S - \Delta S_R &= S(\rho_S') - S(\rho_S) + S(\rho_R') - S(\rho_R) \\ &= S(\rho_S') + S(\rho_R') - S(\rho_S \otimes \rho_R) \\ &= S(\rho_S') + S(\rho_R') - S(\rho_{SR}') \end{aligned}$$

$$S(\rho_S \rho_R') = S(\rho') = I(S':R') \geq 0.$$

$$\Rightarrow \Delta S_S + I(S':R') = -\Delta S_R = S(\rho_R') - S(\rho_R)$$

$$\begin{aligned} \rho_R &= e^{-\beta H} / Z \text{ is thermal state} \\ \Rightarrow S(\rho_R) &= \beta \text{tr}(\rho_R H) + \ln Z \end{aligned}$$



$$\begin{aligned}
&= -\text{tr}(S_R' \ln S_R') - \beta \text{tr}(S_R H) - \ln Z \\
&\quad + \beta \text{tr}(S_R' H) - \beta \text{tr}(S_R' H) \\
&= \beta \Delta Q - \text{tr}\left(S_R' \ln S_R' - S_R' \ln \frac{e^{-\beta H}}{Z}\right) \\
&= \beta \Delta Q - D(S_R' \| S_R). \quad \square
\end{aligned}$$

## 2. The resource theory of nonuniformity

### 2.1. Motivating examples

Standard Landauer erasure:



$$\begin{aligned}
S &= \frac{1}{2} |0\rangle\langle 0| + \frac{1}{2} |1\rangle\langle 1| \\
\downarrow \\
S' &= |0\rangle\langle 0|
\end{aligned}$$

Work cost of bit  
reset:  $\langle W \rangle = k_B T \ln 2$ .

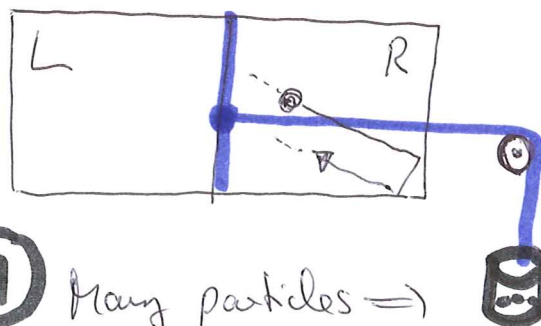
What if  $S$  is any state?



We will show later on: if we have many particles in state  $S$ , the average work cost per particle is  $\langle W \rangle = k_B T S(S)$

von Neumann  
entropy

Complementary picture: the Szilard engine



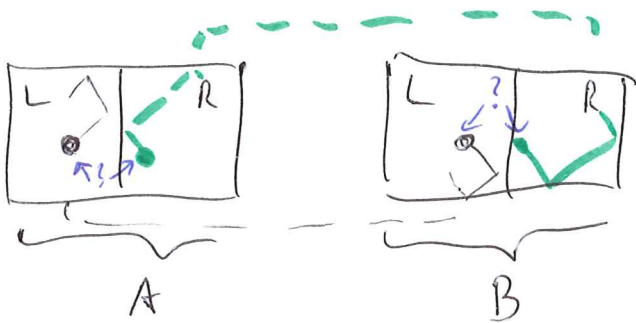
$$S = |R\rangle\langle R|$$

$$\downarrow \\
S' = \frac{1}{2} |L\rangle\langle L| + \frac{1}{2} |R\rangle\langle R|$$

work  
value

④ Many particles  $\Rightarrow$  average work value per particles  $\langle W \rangle = k_B T (1 - S(S))$

Consider two correlated Szilard engines (cf. Bennett 2003)



$$S_{AB} = \frac{1}{2} |LL\rangle\langle LL| + \frac{1}{2} |RR\rangle\langle RR|$$

the local states are  $S_A = \text{Tr}_B S_{AB} = \frac{1}{2} |L\rangle\langle L| + \frac{1}{2} |R\rangle\langle R| = \frac{1}{2} \mathbb{1}$ ,  
 $S_B = \frac{1}{2} \mathbb{1}$ .

$\Rightarrow$  By **acting** locally, we can extract zero work.

But: we can reversibly implement the CNOT operation

$$\begin{aligned} LL &\rightarrow LL \\ LR &\rightarrow LR \\ RL &\rightarrow RR \\ RR &\rightarrow RL \end{aligned}$$

• unitary  $U |RL\rangle = |RR\rangle$   
etc.

$$\begin{aligned} \tilde{S}_{AB} &:= U S_{AB} U^\dagger = \frac{1}{2} |LL\rangle\langle LL| + \frac{1}{2} |RL\rangle\langle RL| \\ &= \left(\frac{1}{2} \mathbb{1}\right) \otimes |L\rangle\langle L| \end{aligned}$$

Now we can extract  $kT \log 2$  of work from engine B!

$\Rightarrow$  Strategy:

1. Extract as many pure (qu)bits as possible (via reversible transformations).
2. Use them for work extraction.

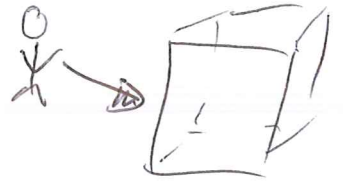
Remarks for step 1:

- $|L\rangle$  and  $|R\rangle$  assumed to have the same energy  $\Rightarrow$  Hamiltonian fully degenerate.
- Laws of nature restrict us to reversible transformations (and this is the only hard restriction).

• We can always introduce extra engines in "thermal" state  $S = \frac{1}{2} |L\rangle\langle L| + \frac{1}{2} |R\rangle\langle R| = \frac{1}{2} \mathbb{I}$  for free

## 2.2. Definition of the resource theory of nonuniformity

Resource theory: an agent has limited control over some physical system



R.T. of nonuniformity: The agent is given some quantum state  $S_{in}$  on a finite-dim. Hilbert space  $\mathcal{H}_{in}$ .

Then he may

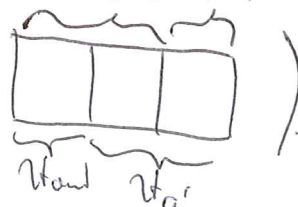
- introduce an extra system  $\mathcal{H}_a$  in the maximally mixed state  $\gamma_a := \frac{\mathbb{I}}{d_a}$ ,  $d_a = \dim \mathcal{H}_a$ , "for free",
- apply arbitrary unitary maps on arbitrary subsystems,  $a, d$
- disregard subsystems, i.e. trace/marginalize over subsystems.

Then every map he can implement by composing these operations can be written in the form

$$S_{out} = \mathcal{E}(S_{in}) = \text{Tr}_{a'} [ U (S_{in} \otimes \gamma_a) U^\dagger ] \quad (*)$$

where  $\mathcal{H}_{in} \otimes \mathcal{H}_a = \mathcal{H}_{out} \otimes \mathcal{H}_{a'}$  and  $U$  is an arbitrary unitary.

(For example with 3 systems:



⑥ Def. Maps of the form (\*) are called noisy operations.



## Remarks:

- Max. mixed states  $\rho$  are free, but "non-uniform" states (e.g. pure states) are valuable resources.  
→  $E(\rho/d_{\text{in}}) = d_{\text{out}}/d_{\text{in}}$  ( $E$  is a unital quantum channel).
- Assumes perfect control (implementation of arbitrary  $U$ )  
→ will explore ultimate limits of what is possible in principle.
- Main question: given  $\rho$  and  $\sigma$ , when does there exist a noisy operator  $E$  with  $E(\rho) = \sigma$ ?  
i.e.  $\rho \xrightarrow{\text{noisy}} \sigma$
- Extraction of pure bits  $|0\rangle \in \mathbb{C}^2$  from  $\rho$ :  
ask for largest integer  $n$  such that  $\rho \xrightarrow{\text{noisy}} |0\rangle\langle 0|^{\otimes n}$   
possibly up to some error  $\epsilon > 0$ .

In the following, we will answer this!

- Later on: introduce Hamiltonians  $H$  to systems, and demand energy conservation  
→ "resource theory of athermality",  
can explicitly compute extractable work without assuming Szilard engines / Landauer's Principle.

1st idea: arbitrary unitaries allow us to diagonalise everything → compute classically!