

Single-shot thermo, lecture 11 ①

Recap last lecture:

Free energy as a transition rate:

$m_n :=$ largest integer with $A^{\otimes n} \xrightarrow{\text{S-Thermal}} B^{\otimes m_n}$, then

$$\lim_{n \rightarrow \infty} \frac{m_n}{n} = \frac{D(S_A \| \gamma)}{D(S_B \| \gamma)} = \frac{F_\beta(S_A) - F_\beta(\gamma)}{F_\beta(S_B) - F_\beta(\gamma)}$$

where $F_\beta(\rho) = \frac{1}{\beta} (\ln Z - S(\rho)) / \beta (= "U - TS")$.

Reduction to classical case is not exactly possible (in contrast to resource theory of nonuniformity):

- not only eigenvalues of state S_S important, but also how eigenbasis of S_S "is rotated with respect to eigenbasis of H_S ".
- Many equivalences of resource theory of nonunif. do not hold in the resource theory of athermality, cf.

$S \rightarrow G$ by Gibbs-preserving map

$\nLeftrightarrow S \rightarrow G$ by a thermal operation.

3.5. From the quantum to the classical case?

Continued.

In light of these problems, let's stick to block-diagonal states, i.e. $[S_S, H_S] = 0 \rightarrow$ everything becomes classical.

Classical system: $A = (p_A, H_A)$

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$p_A \in \mathbb{R}^n$ probability vector, $H_A = (E_1, \dots, E_n) \in \mathbb{R}^n$ energies
A permutation $\pi: \mathbb{R}^n \rightarrow \mathbb{R}^n$ is energy-preserving if
 $\pi[H_A] = H_A$.

Def: A map $D: \mathbb{R}^{d_S} \rightarrow \mathbb{R}^{d_S}$ is a thermal classical operation if there is a system $E = (p_E, H_E)$ with
 $(p_E)_i = \exp(-\beta(H_E)_i) / Z$ and an energy-preserving permutation π on SE with
 $D(p_S) = (\pi[p_S \otimes p_E])_{S'}$.

$P = (p_S, H_S), Q = (q_{S'}, H_{S'})$
Write $P \xrightarrow{\text{thermal}} Q$ if for every $\varepsilon > 0$ there is q_ε with
 $\|q - q_\varepsilon\| < \varepsilon$ and a thermal operation D_ε with
 $D_\varepsilon(p) = q_\varepsilon$. Analogous def. hold for quantum states.

Lemma: Let S, G be quantum states on S .

If S, G are blockdiagonal then

$$S \xrightarrow[\text{quantum}]{\text{thermal}} G \iff \lambda(S) \xrightarrow[\text{classically}]{\text{thermal}} \lambda(G)$$

where $\lambda(S)$ are the eigenvalues of S , ordered in accordance with the ordering of the energy values.

Proof: Maybe in homework.

Theorem (Janzing et al. 2000):

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Let $\mathbf{P} = (p_\bullet, H_\bullet)$ and $\tilde{\mathbf{P}} = (\tilde{p}_\bullet, \tilde{H}_\bullet)$ be classical systems of the same size, i.e. $p_\bullet, \tilde{p}_\bullet \in \mathbb{R}^n$.

Then

$$\mathbf{P} \xrightarrow{\text{thermal}} \tilde{\mathbf{P}}$$

if and only if there exists a stochastic matrix A

$$\text{such that } Ap = \tilde{p} \text{ and } Ag = \tilde{g},$$

where $g = \exp(-\beta H_\bullet) / Z$ and similarly \tilde{g} are the Gibbs states.

Important Theorem! This will be our main tool in the following.

Proof sketch:

" \Rightarrow ": In special case $H_\bullet = \tilde{H}_\bullet$, this means that thermal operations preserve the Gibbs state. Prove analogous to quantum case (cf. homework).

" \Leftarrow ":

Consider special case $n=2$ (i.e. bits).

$A = \begin{pmatrix} a_{00} & a_{01} \\ a_{10} & a_{11} \end{pmatrix}$ stochastic matrix, i.e. $a_{00} + a_{10} = a_{01} + a_{11} = 1$, $a_{ij} \geq 0$. Assume $Ap = \tilde{p}$ and $Ag = \tilde{g}$.

Consider the composite system with Hamiltonian

(4)

$$\underbrace{H + H + H + \dots + H}_{n} + \underbrace{H + \dots + H}_{n} + \underbrace{H}_{\text{output}}$$

input

(in total $(2n+1)$ bits), and a thermal operation

$$D(p) = \left(\Pi \left[p \otimes g^{\otimes n} \otimes \tilde{g}^{\otimes (n+1)} \right] \right)_{2n+2}.$$

← want: this is $\approx \tilde{p}$.

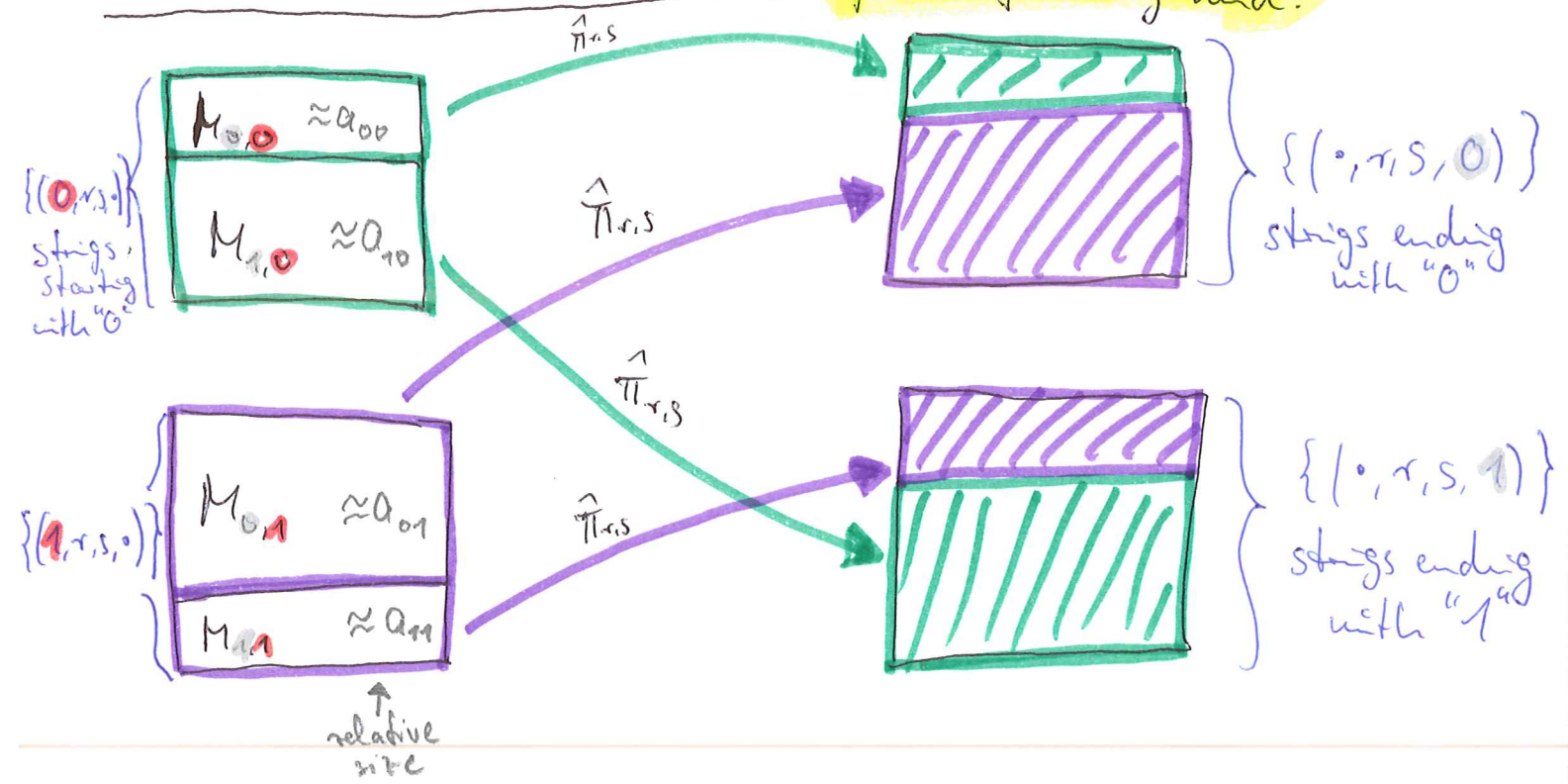
Construction of the energy-preserving permutation Π :

Two binary strings of length $(n+1)$ each are characterized by 4 numbers:

- first bit of the first word, denoted $j \in \{0, 1\}$,
- number of ones in the first word, denoted $r \in [0, n+1]$
- " " " " " second " , denoted $s \in [0, n+1]$
- last bit of the last word, denoted $x \in \{0, 1\}$.

→ 4-tuple (j, r, s, x) .

Fix r, s . Then do a permutation of the following kind:



• Why would this give us what we want?

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If $n \rightarrow \infty$, we can restrict to typical subset.

$$T_n := \{ (j, r, s, x) \mid \frac{r}{n} \approx g_1, \frac{s}{n} \approx \tilde{g}_1 \}$$

• approx. uniform distribution on T_n

\Rightarrow prob. to be in some subset \approx relative size of subset.

\Rightarrow Prob (land in **O-box** after permutation)

$$= \text{Prob}(\text{start in } M_{00} \text{ or } M_{01})$$

$$\approx \text{Prob}(\text{start in } \text{green box}) \cdot a_{00}$$

$$+ \text{Prob}(\text{start in } \text{pink box}) \cdot a_{01}$$

$$\approx p_0 a_{00} + p_1 a_{01} = \sum_j a_{0j} p_j = (A p)_0 = \tilde{p}_0.$$

• Check that we can indeed construct $\tilde{\pi}_{r,s}$:

Left boxes must fit (tightly) into right boxes:

$$\text{Need } \# \{ (\cdot, r, s, 0) \} \approx a_{00} \cdot \# \{ (0, r, s, \cdot) \} + a_{01} \cdot \# \{ (1, r, s, \cdot) \}. \quad (*)$$

$$\text{But } \# \{ (\cdot, r, s, 0) \} \approx \tilde{g}_0 \cdot \# \{ (\cdot, r, s, \cdot) \} \quad \text{(combinatorics for large } n \text{)}$$

$$\text{and } \# \{ (j, r, s, \cdot) \} \approx g_j \cdot \# \{ (\cdot, r, s, \cdot) \}$$

Then (*) follows from $\tilde{g}_0 = a_{00} g_0 + a_{01} g_1$,

$$\text{i.e. } \tilde{g} = A g.$$

(**) Introduce: Consider n bits, $\text{iid} \sim p(1)=p$. (6)

$$\frac{\# \{ \text{strings with } s \text{ ones, last bit zero} \}}{\# \{ \text{strings with } s \text{ ones, last bit arbitrary} \}} = \frac{\binom{n-1}{s}}{\binom{n}{s}}$$

$$\# \{ \text{strings with } s \text{ ones, last bit arbitrary} \} = \binom{n}{s}$$

If $s \approx \alpha \cdot n$, then this tends to $(1-p)$ for $n \rightarrow \infty$.

3.6. d -majorization and d -Lorent curves

Def.: Let $p, p', d, d' \in \mathbb{R}^n$ be probability vectors. We say that (p, d) d -majorizes (p', d') if there is a channel (i.e. stochastic matrix) B with $Bp = p'$ and $Bd = d'$. Then we write: $(p, d) \succ (p', d')$.

Special case $d=d'$, we also write $p \succ_d p'$.

$d=d' = (\frac{1}{n}, \dots, \frac{1}{n})$: standard majorization

Theorem: For probability vectors $p, p', d, d' \in \mathbb{R}^n$, the following conditions are equivalent:

(i) $(p, d) \succ (p', d')$

(ii) for every convex function g ,

$$\sum_i d_i g\left(\frac{p_i}{d_i}\right) \geq \sum_i d'_i g\left(\frac{p'_i}{d'_i}\right),$$

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(iii) the d -Lorenz curve of p is everywhere on or above the d' -Lorenz curve of p' .

(iv) $p \xrightarrow{\text{thermal}} p'$, if initial and final Hamiltonians are such that d and d' are the corresponding Gibbs states.

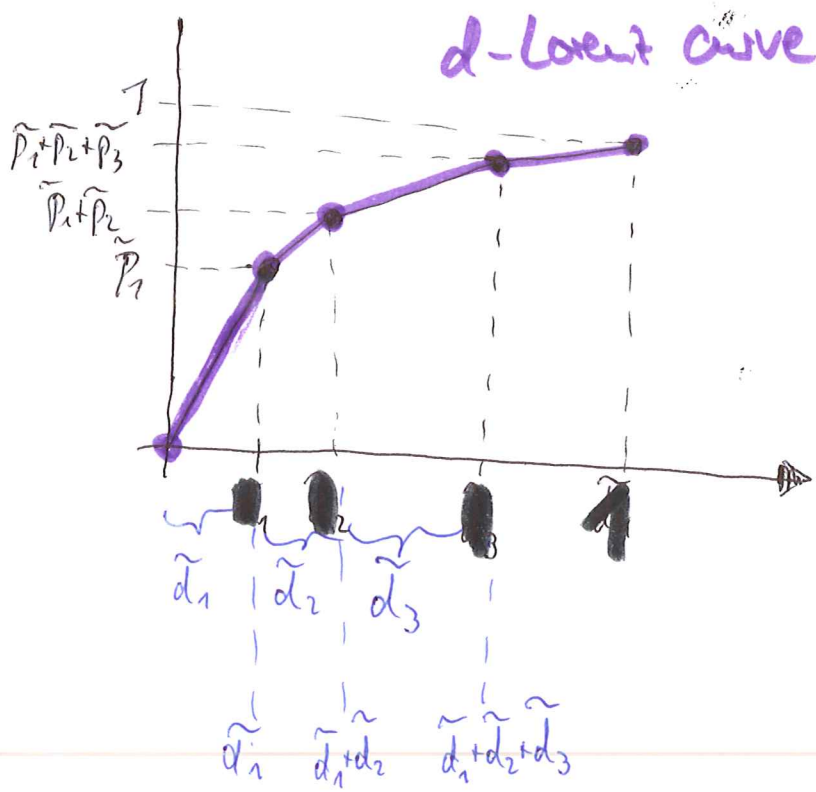
Proof: (i) \Leftrightarrow (iv) is already proven.
 Rest maybe in homework.

d -Lorenz curve (also "Gibbs-rescaled Lorenz curve"):

Consider $\left(\frac{p_i}{d_i}\right)_{i=1, \dots, n}$ and sort them in decreasing order:

$\tilde{p} = \pi(p)$, $\tilde{d} = \pi(d)$ (same permutation) such that

$$\frac{\tilde{p}_1}{\tilde{d}_1} \geq \frac{\tilde{p}_2}{\tilde{d}_2} \geq \frac{\tilde{p}_3}{\tilde{d}_3} \geq \dots \quad (***)$$



Curve is concave due to $(***)$.

Consider condition (ii). For $\alpha > 1$, the function ⑧
 $g(x) := x^\alpha$ is convex, thus

$$\sum_i p_i^\alpha d_i^{1-\alpha} \geq \sum_i (\bar{p}_i)^\alpha (\bar{d}_i)^{1-\alpha}.$$

Def: (Relative Rényi entropy).

For $\alpha \in \mathbb{R} \setminus \{0, 1\}$, set for prob. vectors $p, q \in \mathbb{R}^n$:

$$D_\alpha(p \| q) := \frac{\text{sgn } \alpha}{\alpha - 1} \log \sum_{i=1}^n p_i^\alpha q_i^{1-\alpha}.$$

Furthermore,

$$D_0(p \| q) := \lim_{\alpha \rightarrow 0^+} D_\alpha(p \| q) = -\log \sum_{\substack{i: \\ p_i \neq 0}} q_i$$

$$D_1(p \| q) = \lim_{\alpha \rightarrow 1} D_\alpha(p \| q) = \sum_{i=1}^n p_i (\log p_i - \log q_i)$$

standard relative entropy

$$D_\infty(p \| q) = \lim_{\alpha \rightarrow \infty} D_\alpha(p \| q) = \log \max_i \frac{p_i}{q_i}$$

$$D_{-\infty}(p \| q) = \lim_{\alpha \rightarrow -\infty} D_\alpha(p \| q) = D_\infty(q \| p).$$

Thm: $(p, d) \succ (p', d') \Rightarrow$

$$D_\alpha(p \| d) \geq D_\alpha(p' \| d') \quad \forall \alpha.$$

\Rightarrow Thermal operations decrease all relative-entropy-distances to the Gibbs state.

cf. homework.